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1 recites statutory subject matter. Dependent claims 13-15 are similarly amended.
2 Applicants respectfully request that the §101 rejection of claims 12-15 be
3 withdrawn.

4 Claim 16 has been amended to state that the recited code is “stored in
5 computer readable media and executable on a processor.” This amendment
6 changes the form of claim 16 to the more familiar and acceptable *Beauregard*
7 format. Dependent claims 17 and 18 incorporate these changes by virtue of their
8 dependency on base claim 16. Applicants respectfully request that the §101
9 rejection of claim 16-18 be withdrawn.

10
11 **35 U.S.C. §102**

12 Claims 1 and 4 are rejected under 35 U.S.C. §102 as being anticipated by
13 U.S. Patent 5,835,726 to Shwed et al (Shwed). Applicants respectfully traverse the
14 rejection.

15 The invention concerns a network architecture in which two endpoints
16 communicate via a virtual private network (VPN) on an otherwise public network,
17 such as the Internet, and an intermediary is permitted to inspect the data
18 communication in a secure and trusted manner.

19 In one implementation, the network architecture has an external client and
20 an internal client that exchange encrypted data over a network. The internal client
21 is coupled to the network via a network access point, such as a firewall/proxy
22 server. All three participants have their own pair of public/private keys. An
23 independent key server holds the public keys for all three participants.

24 The external and internal clients establish a virtual private network by
25 negotiating a session key used to encrypt data being exchanged between them.

1 Initially, only the clients know the session key, and not the firewall. To grant the
2 firewall trusted access to the data stream on the VPN, the internal client securely
3 transfers the session key to the firewall. The internal client requests and receives
4 the firewall's public key from the key server and encrypts the session key using the
5 firewall's public key. The internal client then signs the encrypted key by
6 encrypting it using the internal client's private key.

7 The firewall authenticates the signature by decrypting the message using the
8 internal client's public key (obtained from the key server or directly from the
9 internal computer). The firewall then decrypts the session key using its own
10 private key. If the dual decryption yields a valid key, the firewall is assured that
11 the session key was sent by the internal client and was not subsequently altered or
12 tampered with in route.

13 Once the session key is transferred, the firewall is able to decrypt the data
14 stream on the VPN. The firewall can now unintrusively inspect the data stream in
15 a manner that is transparent to the external and internal clients. The claims capture
16 this architecture and new technology.

17 **Claim 1** for example recites a "method for inspecting an encrypted data
18 stream being transferred over a network between two endpoints, the data stream
19 being encrypted using a session key known to both endpoints, the method
20 comprising:

21 securely transferring the session key from one of the endpoints to an
22 intermediary having access to the encrypted data stream;

23 decrypting the encrypted data stream at the intermediary using the session
24 key; and

25 inspecting the data stream following decryption."

1 The method of claim 1 provides for an establishment of a virtual private
2 network (VPN) between two computers (endpoints) where the computers
3 (endpoints) engage in key negotiation process to negotiate a session key (see
4 specification page 9, lines 11-13). With the session key, the endpoints (internal
5 and external clients) are able to encrypt messages and begin an encrypted
6 communication session directly with one another (see specification page 9, lines
7 11-17, Fig. 2). Once the session key is created, one of the endpoints is able to
8 securely share the key with the intermediary to permit trusted inspection.

9 The method of claim 1 is not disclosed by Shwed. Shwed shows host 1 and
10 host 2 computers (also referred to by the Examiner as endpoints) connected to
11 respective private networks. Host 1 and host 2 are secured through respective
12 firewalls. The firewalls connect to one another by way of a public network. See
13 Shwed, col 14, lines 19-39, Fig. 16. Host 1 and host 2 do not directly
14 communicate with one another.

15 Shwed does not disclose that the host computers use a session key known to
16 each other. Instead, Shwed uses firewalls to generate session keys (see Shwed col.
17 15, lines 33-36), not the host computers. The firewalls also perform the
18 encryption. This is pointed out by Examiner that "packet filter modules are
19 installed in firewalls where modification, inspection of packets is performed by
20 encryption of outbound packets and decryption of inbound packets." Shwed at col.
21 13, lines 6-20.

22 Claim 1 recites "using a session key known to both endpoints" and
23 "securely transferring the session key from one of the endpoints to an
24 intermediary." Shwed is silent as to this feature. Since Shwed's system does not
25 allow the host computers to negotiate a session key, there is no way that one of the

1 hosts can securely transfer the session key to an intermediary such as a firewall.
2 Shwed shows a shared session key R used by firewall computers, firewall1 and
3 firewall2, but not between host 1 and host 2 computers. (Shwed at col. 16, lines
4 27-31) Host 1 and host 2 computers in Shwed never see the session key R. In
5 Shwed, session key R is not transferred from *one end point to an intermediary* by
6 encrypting the session key R using a secret common key B. Host 1 and host 2
7 computers are not involved in a key exchange, never see any of the keys, and
8 therefore would not be able to transfer a session key or any key to an intermediary
9 (firewall). For these reasons, claim 1 is patentable over Shwed.

10 The Examiner points out that “Shwed teaches that an endpoint in sending a
11 packet M to a receiving endpoint, first it generates a signature on the packet and
12 then encrypts it using a function key of session key R and finally transmits the
13 encrypted packet over the public network to the receiver endpoint” referring to
14 Shwed col. 18, lines 9-46. As discussed, Shwed shows firewalls encrypting
15 packets, not endpoints. Further, in Shwed only the firewalls see the session key R,
16 the host computers never see session key R.

17 Shwed shows a firewall receiving an encrypted packet and verifies the
18 encrypted packet through decryption of the encrypted packet through decryption of
19 the encrypted packet and verification of the signature. Shwed col. 18, lines 47-67
20 through col. 19, lines 1-3. This session data exchange is “performed by a firewall
21 in receiving an encrypted packet from another firewall.” Shwed col. 18, lines 47-
22 49. The endpoints in Shwed receive “decrypted” packets. Shwed col. 15, lines 5-
23 11. The present invention as recited in claim 1 involves “an encrypted data stream
24 being transferred over a network between two endpoints.” As discussed, endpoints
25 computers and intermediary firewall computers are treated as separate entities.

1 The Examiner in applying Shwed incorrectly treats endpoint computers as being
2 the same as intermediary firewall computers.

3 For these reasons, claim 1 is patentable over Shwed. Applicants
4 respectfully request that the §102 rejection of claim 1 be withdrawn.

5 Dependent claim 4 is allowable by virtue of its dependency on base claim 1.
6 For the reasons given above with respect to claim 1, the systems and methods
7 recited in claim 4 are neither disclosed nor taught by Shwed. Applicants
8 respectfully request that the §102 rejection of claim 4 be withdrawn.

9
10 **35 U.S.C. §103**

11 Claims 2, 3 and 5-20 are rejected under 35 U.S.C. §103 as being
12 unpatentable over Shwed in view of Bruce Schneier, Applied Cryptography,
13 Second Addition, 1996 (Schneier). Applicants respectfully traverse the rejection.

14 Claims 2 and 3 depend from claim 1 and hence incorporate the features of
15 claim 1. As such claims 2 and 3 require “securely transferring the session key
16 from one of the endpoints to an intermediary having access to the encrypted data
17 stream; decrypting the encrypted data stream at the intermediary using the session
18 key; and inspecting the data stream following decryption.”

19 Shwed does not suggest nor teach an architecture where a VPN can be
20 established such that two endpoints can share a session key. The session key in
21 Shwed is known only by and between the intermediary firewalls not the endpoint
22 computers. The endpoints in Shwed do not share a common session key nor are
23 they involved in encryption with one another. Thus, Shwed does not suggest nor
24 teach transferring a session key from one of the endpoints to an intermediary as
25 recited in claim 1. Since the intermediary firewalls perform session key exchange

1 with one another, there would not be a need for an endpoint to transfer the session
2 key to an intermediary. Further, in Shwed, it would be impossible for an endpoint
3 to transfer a session key to an intermediary, since the endpoints never see the
4 session key.

5 Schneier is cited for its teaching of known cryptosystems, in particular key
6 exchange systems. Schneier provides no assistance as to the recited methodology
7 of claims 2 and 3. Accordingly, a combination of Shwed and Schneier fails to
8 teach or suggest the claimed methods. Applicants respectfully request that the
9 §103 rejections of claims 2 and 3 be withdrawn.

10 Claim 5 defines “a method for inspecting an encrypted data stream being
11 transferred over a network between two endpoints and via an intermediary, the
12 data stream being encrypted using a session key known to both endpoints ...
13 passing the signed encrypted session key to the intermediary.” As discussed, the
14 Shwed/Schneier combination does not suggest nor teach encrypted data streams
15 that are transferred between two endpoints using a session key known to the two
16 endpoints. The Shwed/Schneier combination does not suggest nor teach that a
17 session key be passed from an endpoint to an intermediary. Therefore, even in
18 view of Schneier, claim 5 is not obvious. Applicants respectfully request that the
19 §103 rejection of claim 5 be withdrawn.

20 Dependent claim 6 is allowable by virtue of its dependency on base claim 5.
21 Applicants respectfully request that the §103 rejection of claim 6 be withdrawn.

22 Amended claim 7 defines “in a network system having an internal client
23 that exchanges encrypted data with an external client over a network and through a
24 firewall intermediate of the internal and external clients, the encrypted data being
25 encrypted using a session key known to the internal and external clients ... a

1 method executed at the firewall comprising receiving an encrypted and signed
2 session key from the internal client.” As discussed, the Shwed/Schneier
3 combination does not suggest nor teach that an internal client exchange encrypted
4 data with an external client using a session key known to the internal and external
5 clients. The Shwed/Schneier combination further fails to teach that a session key
6 be received by a firewall from an internal client. Applicants respectfully request
7 that the §103 rejection of claim 7 be withdrawn.

8 Dependent claims 8, 9, 10, and 11 are allowable by virtue of their
9 dependency on base claim 7. Applicants respectfully request that the §103
10 rejection of claims 8, 9, 10, and 11 be withdrawn.

11 Claim 12 defines “a network system comprising an internal client and an
12 external client configured to communicate encrypted data over a network ..., the
13 data being encrypted using a session key, the internal client being configured to
14 securely transfer the session key to the intermediary.” As discussed, the
15 Shwed/Schneier combination does not suggest nor teach that an internal client
16 exchange encrypted data with an external client using a session key known to the
17 internal and external clients. The Shwed/Schneier combination further fails to
18 teach that the internal client be configured to transfer such a session key to an
19 intermediary. Applicants respectfully request that the §103 rejection of claim 12
20 be withdrawn.

21 Dependent claims 13, 14, and 15 are allowable by virtue of their
22 dependency on base claim 12. Applicants respectfully request that the §103
23 rejection of claims 13, 14, and 15 be withdrawn.

24 Claim 16 defines “a software architecture for a network system having two
25 endpoints that exchange encrypted data over a network and through an

1 intermediary, the encrypted data being encrypted using a session key known to the
2 endpoints comprising: endpoint-resident code stored on computer readable media
3 and executable on a processor to encrypt the session key using a public key from a
4 public/private key pair associated with the intermediary and to sign the encrypted
5 session key with a digital signature, the endpoint-resident code being capable of
6 sending the signed and encrypted session key to the intermediary; and
7 intermediary-resident code stored on computer readable media and executable on
8 the processor to authenticate the digital signature and decrypt the encrypted session
9 key using a private key from the public/private key pair associated with the
10 intermediary, the intermediary-resident code using the session key to decrypt the
11 encrypted data as it is being exchanged between the two endpoints.” As
12 discussed, the Shwed/Schneier combination does not suggest nor teach that an
13 internal client exchange encrypted data with an external client using a session key
14 known to the internal and external clients. The Shwed/Schneier combination fails
15 to teach that a session key be known to the endpoints. The Shwed/Schneier
16 combination fails to teach endpoint-resident code being capable of sending the
17 signed and encrypted session key to the intermediary. The Shwed/Schneier
18 combination further fails to teach intermediary-resident code using the session key
19 to decrypt the encrypted data as it is being exchanged between the two endpoints.
20 Applicants respectfully request that the §103 rejection of claim 16 be withdrawn.

21 Dependent claims 17 and 18 are allowable by virtue of their dependency on
22 base claim 16. Applicants respectfully request that the §103 rejection of claims 17
23 and 18 be withdrawn.

24 Claim 19 defines “a network system having an internal client that
25 exchanges encrypted data with an external client over a network and through a

1 firewall intermediate of the internal and external clients, the encrypted data being
2 encrypted using a session key known to the internal and external clients ... passing
3 the signed and encrypted session key to the intermediary.” As discussed, the
4 Shwed/Schneier combination does not suggest nor teach that an internal client
5 exchange encrypted data with an external client using a session key known to the
6 internal and external clients. The Shwed/Schneier combination further fails to
7 teach the session key being passed to the intermediary. Applicants respectfully
8 request that the §103 rejection of claim 19 be withdrawn.

9 Claim 20 defines “a network system in which an encrypted data stream is
10 transferred over a network between two endpoints and via an intermediary, the
11 data stream being encrypted using a session key known to both endpoints
12 ...securely transferring the session key from one of the endpoints to an
13 intermediary.” As discussed, the Shwed/Schneier combination does not suggest
14 nor teach that an internal client exchange encrypted data with an external client
15 using a session key known to the internal and external clients. The
16 Shwed/Schneier combination further fails to teach that the session key be
17 transferred from one of the endpoints to an intermediary. Applicants respectfully
18 request that the §103 rejection of claim 20 be withdrawn.

1
2 **CONCLUSION**

3 All pending claims 1-20 are in condition for allowance. Applicant
4 respectfully requests reconsideration and prompt issuance of the subject
5 application. If any issues remain that prevent issuance of this application, the
6 Examiner is urged to contact the undersigned attorney before issuing a subsequent
7 Action.

8
9 Respectfully Submitted,

10
11 Dated: 10/29/02

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MARKED UP VERSION OF PENDING CLAIMS UNDER 37 C.F.R. §

1.121(C)(1)(ii):

Amend claim 3, 7, 12-19 as follows and in accordance with 37 C.F.R. § 1.121(c)(1)(ii), by which the Applicant submits the following marked up version only for claims being changed by the current amendment, wherein the markings are shown by brackets (for deleted matter) and/or underlining (for added matter):

3. (Amended Once) A method as recited in claim 1, wherein securely transferring comprises:

encrypting the session key using a public key associated with the intermediary;

signing the encrypted session key using a private key associated with the [intermediary]one of the endpoints; and

sending the signed and encrypted session key to the intermediary.

7. (Once Amended) In a network system having an [external]internal client that exchanges encrypted data with an external client over a network and through a firewall intermediate of the internal and external clients, the encrypted data being encrypted using a session key known to the internal and external clients, a method executed at the firewall comprising:

receiving an encrypted and signed session key from the internal client, the encrypted and signed session key bearing a digital signature of the internal client;

authenticating the digital signature as belonging to the internal client;

decrypting the session key; and

1 decrypting the encrypted data being exchanged between the internal and
2 external clients using the session key.

3
4 12. (Amended Once) A network system comprising:
5 an internal client device and an external client device configured to
6 communicate encrypted data over a network using virtual private network
7 communication, the data being encrypted using a session key;
8 an intermediary device having access to the encrypted data being
9 communicated between the internal client device and the external client device;
10 the internal client device being configured to securely transfer the session
11 key to the intermediary device; and
12 the intermediary device being configured to decrypt the data using the
13 session key and to inspect the data.

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15 13. (Amended Once) A network system as recited in claim 12, wherein
16 the internal client device encrypts the session key prior to sending it to the
17 intermediary device.

18
19 14. (Amended Once) A network system as recited in claim 12, wherein
20 the internal client device encrypts and signs the session key prior to sending it to
21 the intermediary device.

22
23 15. (Amended Once) A network system as recited in claim 12, wherein
24 the intermediary device stores the data in unencrypted form.
25

1 16. (Amended Once) A software architecture for a network system
2 having two endpoints that exchange encrypted data over a network and through an
3 intermediary, the encrypted data being encrypted using a session key known to the
4 endpoints, comprising:

5 endpoint-resident code stored on computer readable media and executable
6 on a processor to encrypt the session key using a public key from a public/private
7 key pair associated with the intermediary and to sign the encrypted session key
8 with a digital signature, the endpoint-resident code being capable of sending the
9 signed and encrypted session key to the intermediary; and

10 intermediary-resident code stored on computer readable media and
11 executable on the processor to authenticate the digital signature and decrypt the
12 encrypted session key using a private key from the public/private key pair
13 associated with the intermediary, the intermediary-resident code using the session
14 key to decrypt the encrypted data as it is being exchanged between the two
15 endpoints.
16

17 17. (Amended Once) A software architecture as recited in claim 16,
18 wherein the intermediary-resident code inspects the data in unencrypted form.
19

20 18. (Amended Once) A software architecture as recited in claim 16,
21 wherein the intermediary-resident code stores the data in unencrypted form.
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1 19. (Amended Once) In a network system having an [external]internal
2 client that exchanges encrypted data with an external client over a network and
3 through a firewall intermediate of the internal and external clients, the encrypted
4 data being encrypted using a session key known to the internal and external clients,
5 computer-readable media distributed at the internal client and the firewall storing
6 computer-executable instructions for:

7 encrypting the session key at the internal client;

8 signing the encrypted session key with a digital signature associated with
9 the internal client;

10 passing the signed and encrypted session key to the intermediary;

11 authenticating, at the intermediary, the digital signature of the internal
12 client;

13 decrypting the session key at the intermediary;

14 decrypting, at the intermediary, the encrypted data using the session key;

15 and

16 inspecting the data in route between the internal and external clients.
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1 MARKED UP VERSION OF SPECIFICATION UNDER 37 C.F.R. §

2 1.121(B)(1)(iii):

3
4 Page 10, second paragraph

5 Fig. 3 shows an exemplary implementation of a computer, such as the
6 external client 42, the internal client 44, firewall [58]48, or key server 66. The
7 host computer is a general-purpose computing device in the form of a conventional
8 personal computer 100 that is configured to operate as a network server (in the
9 case of the firewall and key server computers) or as a client computer (in the case
10 of the external and internal clients).
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